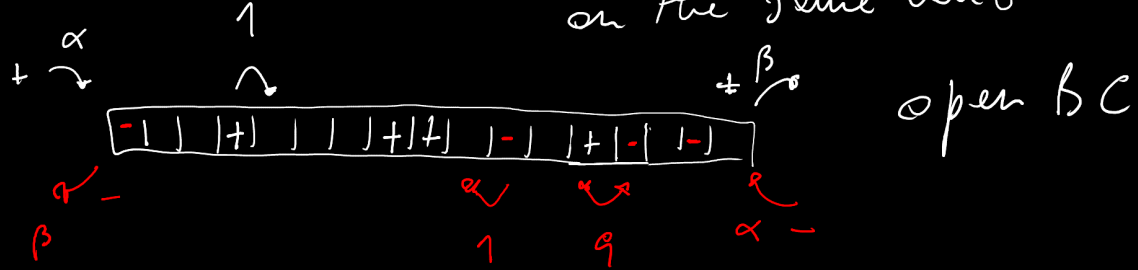


BRIDGE MODEL \rightarrow TASEP with 2 particle types (+, -) moving in opposite directions on the same lane



Symmetry upon exchange + with -

if $(\beta \ll 1 ; \alpha, \rho \sim 1)$: both + and - should stay in the HD phase

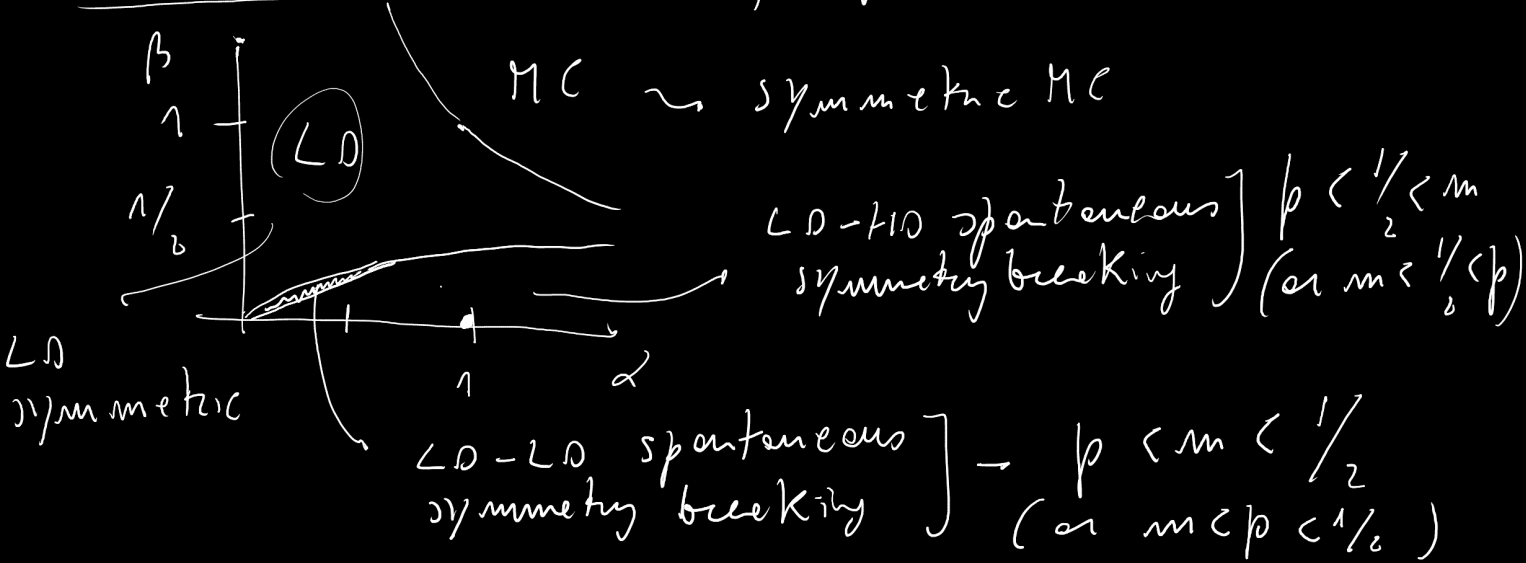
$$p + m \leq 1$$

densities of positive (p) and negative (m) particles

whereas $\underline{p_{HD} + m_{HD} > 1}$

Spontaneous symmetry breaking

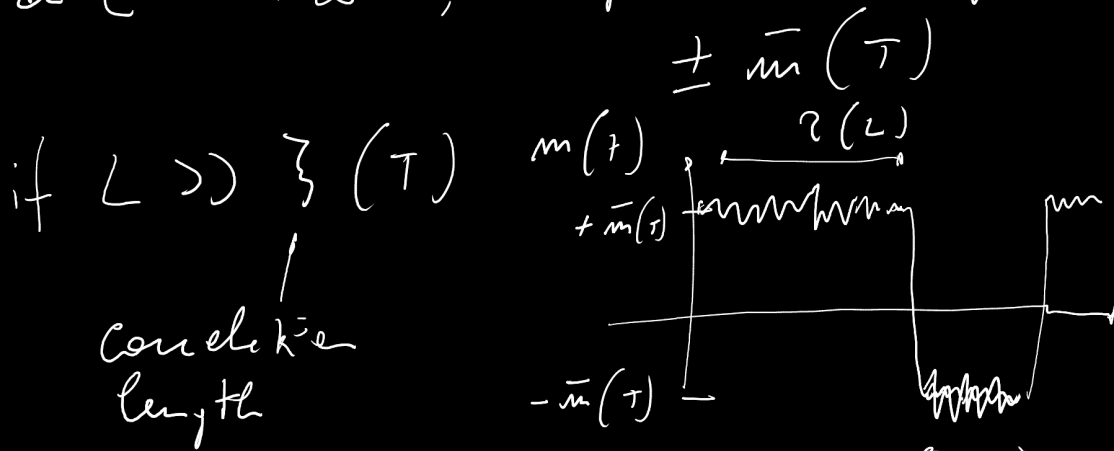
PHASE DIAGRAM (in α, β plane)



SSB proved exactly for $\beta \ll 1 ; \alpha = \rho = 1$

How can SSB be proved? Ising model in $d=2$
 finite size L at $T < T_c$ $h=0$

for $L \rightarrow +\infty$, expected values for magnetization



$\langle \xi(L) \rangle =$ characteristic time for magnetization inversion

$\langle \xi(L) \rangle \sim \exp(L)$ \Rightarrow Spontaneous symmetry breaking!

Ising model: switch on a symmetry breaking field ($h > 0$) \Rightarrow NO COEXISTENCE ANYMORE

Bridge model: switch on a field small enough \Rightarrow COEXISTENCE

INFORMATION, PHYSICS & COMPUTATION

Entropy \rightarrow Information \rightarrow Data compression / Data transmission

X stochastic discrete variables assuming values $x \in X$ (finite discrete set)

$\{p(x)\}_{x \in X}$ PDF $p: X \rightarrow [0, 1]$

$\sum_{x \in X} p(x) = 1$ expectation $E f = \sum_{x \in X} p(x) f(x)$
value for any real-valued function $f(X)$

$$\text{Var } f(X) = E \{ f^2(X) \} - \{ E f(X) \}^2$$

$A \subseteq X$ (subset) $\int p(x) dx$

$$P(X \in A) = \int_{x \in A} dp(x) = \int \mathbb{I}(x \in A) dp(x)$$

indicator function $\mathbb{I}(s) = \begin{cases} 1 & \text{statement } s \text{ true} \\ 0 & \text{statement } s \text{ false} \end{cases}$

Gaussian PDF $p(x) = \frac{1}{\sqrt{2\pi\sigma^2}} \exp\left(-\frac{(x-m)^2}{2\sigma^2}\right)$

$$E X = m ; \text{ var } X = E (x-m)^2 = \sigma^2$$

Jensen's inequality $X \subseteq \mathbb{R}$; f convex function

$$\left[\forall x, y \quad \forall \alpha \in [0, 1]: f(\alpha x + (1-\alpha)y) \leq \alpha f(x) + (1-\alpha)f(y) \right]$$

$$\Rightarrow E f(X) \geq f(E X)$$

SHANNON ENTROPY

H_X discrete random variable X with PDF $p(x)$

$$H_X \equiv - \sum_{x \in X} p(x) \log_2 p(x) = \mathbb{E} \left[\log_2 \left(\frac{1}{p(x)} \right) \right]$$

$\log_2 \rightarrow$ Shannon entropy \rightarrow bits

Entropy: measure of uncertainty \rightarrow missing information

$H_X \approx \log_2$ (# of different values the variable may take)

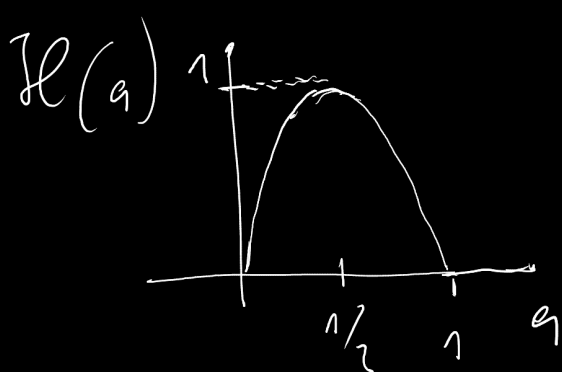
Fair coin $\rightarrow H = 1 = \log_2 2$ ($N = 2$)

M fair coins $\rightarrow H = \log_2 N = M$ (M bits!)
 ($N = 2^M$) (entropy is additive)

fair dice
 (N faces) $\rightarrow H = \log_2 N$

"Unfair" coin (Bernoulli process): $p(0) = q$
 ($q \neq 1/2$) $p(1) = 1 - q$

$$H_X = -q \log_2 q - (1-q) \log_2 (1-q) \equiv \mathcal{H}(q)$$



$q = 1/2 \rightarrow \mathcal{H} = 1$ (fair coin)

$q = 0, 1 \rightarrow \mathcal{H} = 0$ fixed result

Kullback-Leibler (KL) divergence defined between two PDFs $p(x), q(x)$ over the same \mathcal{X}

$$D(q \parallel p) = \sum_{x \in \mathcal{X}} q(x) \log_2 \left[\frac{q(x)}{p(x)} \right]$$

(i) $D(q \parallel p)$ is convex in q $\left(\begin{array}{l} 0 \log 0 = 0 \\ 0 \log \left(\frac{0}{0} \right) = 0 \end{array} \right)$

(ii) $D(q \parallel p) \geq 0$

(iii) $D(q \parallel p) = 0 \iff q(x) = p(x)$

↳ Jensen's with $-\log(x)$ is convex

$$-D(q \parallel p) = \mathbb{E}_q \log_2 \left[\frac{p(x)}{q(x)} \right] \leq \underbrace{\log_2 \mathbb{E}_q \left[\frac{p(x)}{q(x)} \right]}_{\geq 1}$$

$D(q \parallel p) \sim$ "distance" between q, p

$D(q \parallel p) \neq D(p \parallel q)$ (asymmetric)

Go back to Shannon entropy $H_x = - \sum_x p(x) \log_2 p(x)$

(1) $H_x \geq 0$ (2) $H_x = 0 \iff \mathcal{X}$ is certain

$$(p(x) = \delta_{x, x_0})$$

(3) H_x is maximum when all N events are equiprobable

$$\bar{p}(x) = \frac{1}{N}$$

$$\rightarrow H_x(\bar{p}) = \log_2 2^N \text{ generic } p(x)$$

$$D(p \parallel \bar{p}) = \sum_x p(x) \log_2 [N p(x)]$$

$$= \log_2 N - H(p) \geq 0$$

$$\Rightarrow H(p) \leq \log_2 N = H(\bar{p})$$

④ X, Y independent random variables

$$p_{x,y}(x, y) = p_x(x) p_y(y)$$

$$H_{X,Y} = - \sum_{x,y} p(x) p(y) \log_2 [p(x) p(y)]$$

$$= H_x + H_y \quad (\text{additivity rule for independent events})$$

⑤ any pair X, Y is general correlated

$$H_{X,Y} \leq H_x + H_y \quad \left(\text{proven from } D(p_{X,Y} \parallel p_x p_y) \geq 0 \right)$$

correlation \Rightarrow less entropy

\Rightarrow more 'a priori' info on X, Y

⑥ Additivity rule for composite events

$$X = X_1 \cup X_2; \quad X_1 \cap X_2 = \emptyset$$

$$q_1 = \sum_{x \in X_1} p(x) \quad (\text{prob. } X_1) \quad \parallel \quad q_1 + q_2 = 1$$

$$q_2 = \sum_{x \in X_2} p(x) \quad (\text{prob. } X_2)$$

$$\tau_1(x) = p(x)/q_1 \quad \left(\begin{array}{l} \text{conditional prob. for } x \\ \text{give } x \in X_1 \end{array} \right)$$

$$\tau_2(x) = p(x)/q_2 \quad \left(\begin{array}{l} \text{" " " for } x \\ \text{give } x \in X_2 \end{array} \right)$$

$$H_x = \mathcal{H}(q_1) + \tilde{H}(q, \tau)$$

$$\mathcal{H}(q_1) = -q_1 \log q_1 - q_2 \log q_2 \quad \left\{ \begin{array}{l} \text{entropy} \\ \text{related} \\ \text{to subset} \\ \text{choice} \end{array} \right.$$

$$\tilde{H}(q, \tau) = -q_1 \sum_{x \in X_1} \tau_1(x) \log_c \tau_1(x) - q_2 \sum_{x \in X_2} \tau_2(x) \log_c \tau_2(x)$$

entropy related to event choice within subsets

generalized by chain rule for conditional entropies